HW Assignment
Copy Collection

1) Number using original algorithm:

For each root r: copy(r)

1 Copy A object to A’ memory location
16 Copy B object to B’ memory location
3 Copy C object to C’ memory location
5 Copy D object to D’ memory location
11 Copy E object to E’ memory location

2 Update As header with A’ address
17 Update Bs header with B’ address
4 Update Cs header with C’ address
6 Update Ds header with D’ address
12 Update Es header with E’ address

15 Return A’ address and update obj1 val
18 Return B’ address and update obj1 val
9 Return C’ address
7 Return D’ address
13 Return E’ address

10 Update A’ left field with C’s address
14 Update A’ right field with E’s address
8 Update C’ left field with D’s address
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For each root \( r \): \( \text{copy}(r) \)

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16. Copy B object to B’ memory location
3. Copy C object to C’ memory location
5. Copy D object to D’ memory location
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2. Update As header with A’ address
17. Update Bs header with B’ address
4. Update Cs header with C’ address
6. Update Ds header with D’ address
12. Update Es header with E’ address

15. Return A’ address and update \( \text{obj1} \) val
20. Return B’ address and update \( \text{obj2} \) val
9. Return C’ address
7, 18. Return D’ address
13. Return E’ address

10. Update A’ left field with C’s address
14. Update A’ right field with E’s address
8. Update C’ left field with D’s address

19. Update B’ field with D’s address

Here’s a case that is a bit different – B has a field that refers to D
Copying Collector pros/cons

• Strengths
  ■ No fragmentation
    ‣ Fast allocation (bump pointer)
  ■ Handles variable-sized objects naturally
  ■ No overhead on pointer updates (operations)
  ■ Multiple options for placing objects during collection

• Weaknesses
  ■ Required address space is doubled compared with non-copying collectors
    ‣ Performance degrades as residency increases (twice as quickly as mark&sweep b/c half the space)
  ■ Touches every page (VM) of heap regardless of user program residency
    ‣ Unless both semispaces can be held in memory simultaneously
    ‣ Disrupts program locality (memory hierarchy and OS paging)
  ■ Pause for stop-the-world: even longer b/c of copying
  ■ Additional space in object header for forwarding pointer