Underlying computer system = hardware + software

Thanks to Chandra Krintz and Kevin Sanft, for this figure and some other parts of these lecture notes.

Devices alert the CPU that they have a request (keyboard input, disk read return) via an interrupt (signal sent on the bus). This interrupts what the CPU is doing so that it can respond to (execute code for) the request (interrupt handler).
Program instructions and data are in memory
  - CPU tracks which instruction it’s on using a dedicated register (PC) which holds the address of the instruction

CPU stores the next few instructions in a cache – much faster to access than memory
  - Similarly stores data used by the instructions in a data cache
  - For even faster access, the CPU stores some data values and addresses in registers (fewer in number than cache entries and even faster to access than cache)

CPU components (hardware registers, ALU, bus) all use same data width (e.g., 32 bit or 64 bit)
• System bus = address bus + data bus + other signals (wires)
  – CPU requests the next instruction address by putting it on the address bus (wires connected to pins)
  – CPU requests data used by the instruction (operands) by putting the addresses on the data bus
• CPU toggles other pins to identify which devices (memory, IO) it wishes to access – and whether it wants to read or write
• Devices use special wires/pins to alert the CPU that the data that the CPU requested are ready
  – The CPU doesn’t block after a request, it goes onto another task until the device “interrupts” it with the data.
Things to ponder

- How are all of these computer operations managed effectively?
  - After all, the CPU just responds to the next instruction. So how are all the instructions managed, especially when there are many clients (users, processes)?

- How are we – and our simple programs – able to deal with such a complex system?
  - Don’t we need an intermediary?
Operating systems: two views

- **Top-down view:** an OS is software that isolates us from the complications of hardware resources
  - In other words, an OS is an application programmer’s and a user’s *interface* to computer operations

- **Bottom-up view:** an OS is software that allocates and de-allocates computer resources – efficiently, fairly, orderly and securely
Types of operating systems

- Single-user, single-process – i.e., one customer, and one job at a time
- Single-user, multi-process – one workstation, but lots of stuff running
  - Actually the CPU handles just one process at any moment – jobs are swapped in/out in “time slices”
- Multi-user, multi-process – e.g., Unix/Linux
  - Same idea, but much more swapping to do
  - And added fairness, efficiency and security concerns
Unix history (Linux prequel)

- AT&T Bell Labs – **System V** standard
  - 1969-70: Ken Thompson wrote Unix in “B”
  - 1972: Dennis Ritchie developed C – a better B
  - Unix rewritten in C, 1973
  - … eventually System V, 1983

- UC Berkeley – **BSD** standard
  - Started with a copy of System IV, late 1970s
  - Lots of changes/additions in 1980s
  - Now FreeBSD

- Open source – **Linux**, since early 1990s
Unix philosophy (same as C)

- **Small is beautiful**
  - Each program does just one thing
  - Pipe commands (or use successive functions in C) to accomplish more complicated things
  - Less typing is best (using 1970s computers)
    - That’s why so many commands are short (ls, cp, mv, …)

- **Users/programmers know what they are doing**
  - That’s what makes the brevity sufficient
  - Means very few restrictions (or safety nets) apply
Linux

- Linus Torvalds created it as a Finnish undergraduate student
- Posted on Internet in 1991
  - Open source – licensed under GPL
  - Version 1.0 in 1994; version 2.2 in 1999
  - 1000’s of programmers working to enhance it
- When programmers worldwide can read, modify, and redistribute a program’s source code, it evolves.
  - People improve it, adapt it, fix bugs, …
What is Linux?

- A fully-networked Unix-like operating system
- Multi-user, multitasking, multiprocessor system
  - Fundamental in the system’s design and implementation
- Has both command-line and graphical interfaces
- Coexists with other operating systems
- Runs on multiple platforms
- Distribution includes the source code
- Can download it free from the Internet!
The Linux System

User commands includes executable programs and scripts

The shell interprets user commands. It is responsible for finding the commands and starting their execution. Several different shells are available. “Bash” is popular and what we will use.

The kernel manages the hardware resources for the rest of the system

Set of data structures (usually on a disk) that holds directories of files. All devices are accessed like they are files on disk (open/close, read/write).

<table>
<thead>
<tr>
<th>User commands</th>
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<tbody>
<tr>
<td>Shell</td>
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<td>Kernel</td>
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<td>File systems</td>
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</tbody>
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Software that makes use of all of the functionality that each device provides. Drivers implement the file interface (open/close, read/write) so that processes can access the device(s). One driver can support 1+ similar devices.

Thanks again to Chandra Krintz and Kevin Sanft.
Linux kernel – the actual OS

- Manages processes
  - Starts, stops, suspends, swaps, manages inter-process communication, …
  - Maintains their state
- Manages files (and directories)
- Manages main memory
- Manages disk operations
- Delegates to CPU(s), printers, other I/O devices
CPU scheduling

- Kernel sends interrupt to a process to give another process a turn to use the CPU.
- Processes can give up CPU when they don’t need it (e.g. waiting on I/O device).
Processes request kernel services

- Using system calls (read, write, fork, …)
  - OOP idea: these are the kernel’s interface
  - Processes access devices just like files – that’s how they are represented by the kernel, and they occupy places in the file system
    - Use open, close, read, write, release, seek, …

- Or indirectly, by using shell commands or libraries/programs that use system calls
Linux file system

- Rooted, hierarchical
  - Data files are stored in directories
- A file’s (full) pathname starts at the root
  - /etc/passwd
  - /home/neale/b
Some “big picture” ideas

A *simple* computer model

This and the next six figures derived from B. Molay’s *Understanding Unix/Linux Programming*, Pearson 2003.
An example program

```c
#include <stdio.h>
int main(void) {
    int c;
    while ( (c = getchar()) != EOF )
        putchar(c);
}
```
More realistic computer model
How connected? Not like this!
OS manages everything!
OOP idea: OS provides services
User interface is the **shell**

![Diagram showing user space, shell, kernel, disk, and user interaction](image-url)
Shell

- A program that runs in a terminal and provides a command-line interface for user
- An interpreter that executes user commands
- Also a powerful programming language
  - Shell script – a sequence of commands in a file
- Lots of different shells to choose from
  - sh, csh, tcsh, bash …
  - We’ll focus on bash (and sh scripts) in this course
Special file names

- . (by itself) The current directory
  - .a is the same as a
- .. The parent (toward root) directory
  - ../jane/x go up one level then look in directory named jane for x
- ~ Your home directory
  - ~harvey Username harvey’s home directory
- Have to “escape” spaces with a backslash
  - my\ file\ name\ with\ spaces
  - Moral: don’t use spaces in file or directory names!
Object-oriented perspective

Operating system = computer interface

Shell/libraries/system calls = OS interface

Will return to OS topics (processes, ...) in upcoming lectures. Now: OO intro.
Objects

- Include *things*
  - Stack, queue, list, …
  - Window, spaceship, recipe, …
- Also include *concepts*
  - Power, trajectory, mood, …
- Can represent people, places, roles, …
- In programming: an object is a software entity *encapsulating* data and/or methods
Imperative programming (not OOP)

- Data, and the operations that manage the data are separate entities (physically and logically)

- What are implications of this programming style?
Kay’s Description of OOP

1. Everything is an object.
2. Objects perform computations by making requests of each other through the passing of messages.
3. Every object has its own memory, which consists of other objects.
4. Every object is an instance of a class. A class groups similar objects.
5. The class is the repository for behavior associated with an object.
6. Classes are organized into a singly-rooted tree structure, called an inheritance hierarchy.

Alan Kay: “Simple things should be simple, complex things should be possible.”
Solving problems *with* objects

- First decide what objects are needed
  - Instead of what functions are required
  - And instead of how specifically to handle data
- Then give each object **responsibilities**
  - Which probably include storing some data and performing some functions
- Finally, have objects interact by sending **messages** (usually method calls) to one another
  - i.e., they collaborate to fulfill responsibilities
Budd’s “real life” example

- Budd decides to send flowers to his grandmother
- First he selects an agent: Flo, a capable florist
  - Then he sends a message to Flo – not unlike:
    ```
    flo.sendBouquet(1, &grandma);
    ```
- The next step is Flo’s responsibility
  - Budd does not participate in this part of the process
  - Likely that many other agents do participate though!
- Finally Flo probably sends a message to Budd:
  ```
  budd.pay(bouquetPrice, this);
  ```
Elements of OOP - Objects

- 1. Everything is an object
  - Actions in OOP are performed by agents, called *instances* or *objects*.

- Several agents in the example scenario, including Budd, Grandma, Flo, the florist in Grandma’s city, driver, flower arranger, grower
  - Each agent has a part to play, and the result is produced when all work together in the solution of a problem.
2. Objects perform computations by making requests of each other through the passing of messages.
   - Actions in OOP are produced in response to requests for actions, called messages. An instance may accept a message, and in return will perform an action and return a value.

To begin the process of sending the flowers, Budd gives a message to Flo. She in turn gives a message to the florist in Grandma’s city, who gives another message to the driver, and so on.
Information hiding

- Notice how a user of a service being provided by an object, need only know the name of the messages that the object will accept.
  - They need not have any idea how the actions performed in response to these requests will be carried out.

- Having accepted a message, an object is responsible for carrying it out.
Receivers and behavior

- Messages differ from traditional function calls in two very important respects:
  a) A designated receiver accepts the message
  b) The interpretation of the message may be different, depending upon the receiver

- Although different objects may accept the same message, the actions (behavior) the object will perform will likely be different
  - Might not even know what behavior to perform until run-time – a form of late binding
Elements of OOP – Recursive Design

- 3. Every object has its own memory, which consists of other objects.
  - The structure of the part mirrors the structure of the larger unit.

- Principle of non-interference: “Ask not what you can do to your data structures, but ask what your data structures can do for you.” (Budd)
4. Every object is an instance of a class. A class groups similar objects.
   - Flo is an *instance* of the *class* Florist

5. The class is the repository for behavior associated with an object.
   - All objects that are instances of a class use the same method in response to similar messages.
Elements of OOP - Inheritance

- 6. Classes are organized into a *singly-rooted* tree structure, called an inheritance hierarchy.

- Data and *general* behavior at one abstraction level extend to lower levels.
  - But can *override* behavior (a later topic).