

Correlated And Time-Varying Events

- * So far, events and statistics are
 - □ independent
 - □ time (space) invariant
- Many applications violate the above
 speech signals
 images (textures)
- Need a more sophisticated model



Possibilities

This problem (and its variants) has been studied in many fields over a long period of time

We will briefly discuss

- Finite automata (computer science, syntactic pattern recognition)
- □ Markov chain (mathematics, statistics)
- Hidden Markov model (signal processing, speech processing)



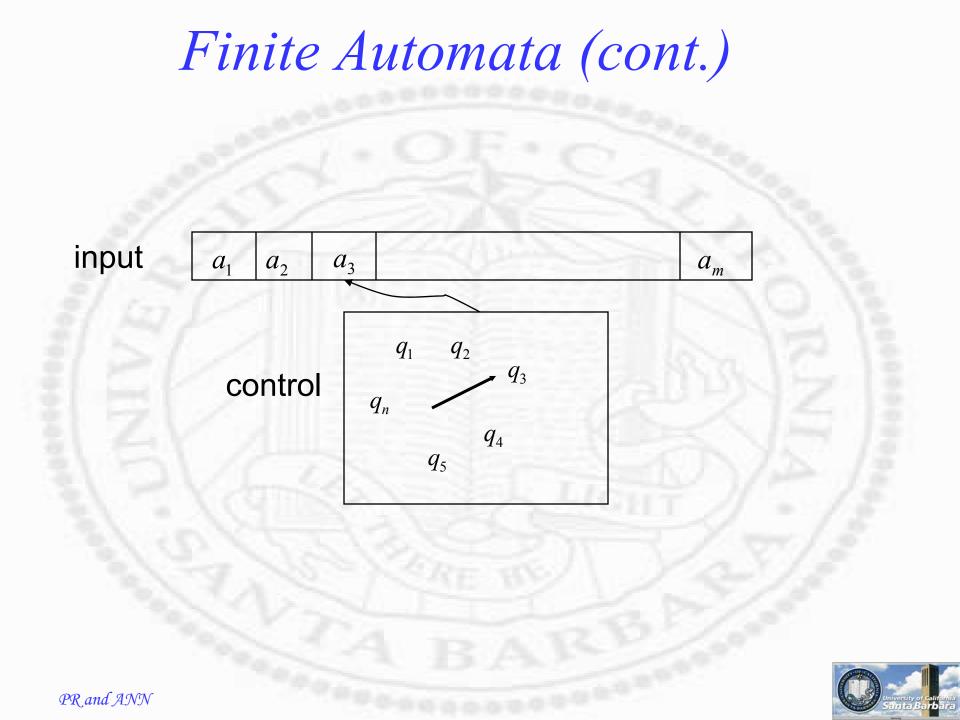
Finite Automata

As a language (or more generally, a *pattern*)
 recognition device

A machine with a *fixed* number of states (memory)
It parses a string (a *pattern*) and based on the current state and current input to decide on the next state

Seventually, when the string (the *pattern*) is exhausted, the machine will halt and either accept (recognize) or reject (not recognize) the input





Finite Automata (cont.)

* A deterministic finite automata is a quintuple $M = (K, \Sigma, \delta, s, F)$

Kis a finite set of states Σ is an alphabet $s \in K$ is the initial state $F \subset K$ is the set of final states $\delta(K \times \Sigma) \rightarrow K$ transition function



Example

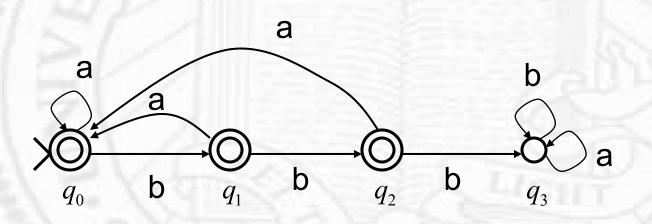
A finite automata that accepts all strings in {a,b}* that have an even number of b's
alphabet: {a,b}
state: {q₀, q₁}
initial state: q₀
final states: {q₀}
transition rules

q	σ	$\delta(q,\sigma)$
q_0	а	q_{0}
q_0	b	q_1
q_1	а	q_1
q_1	b	q_{0}



More Example

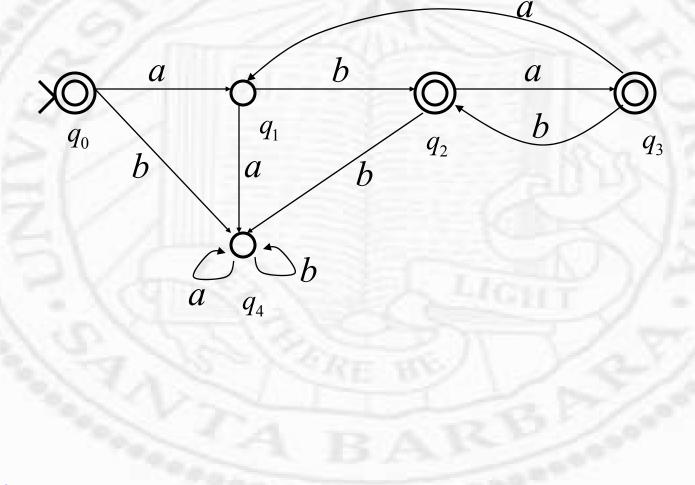
A finite automata that accepts strings over {a,b}* that does not contain three consecutive b's





More Example

* A finite automata that accepts $L = (ab \cup aba)^*$





An Opposite Definition

- Grammar: as a language (or pattern)
 generation device
- ★ A grammar is a quadruple
 □ V an alphabet
 □ Σ subset of V (terminals)
 □ R : (V Σ) → V^{*}rewriting rules
 □ S start symbol



An Opposite Definition

✤ A language

all the strings (patterns) that can be generated by applying the rewriting rules from the start symbol

$L(G) = \{ w \mid w \in \Sigma^*, S \Longrightarrow w \}$



Regular Grammar

The rewriting rule must be

$$R: (V - \Sigma) \to \Sigma^*((V - \Sigma) \cup \{e\})$$

At most one non-terminal on the right-hand side Non-terminal, if present, must be the last symbol in the string



Example

 $a(a^* \cup b^*)b$

R: $S \rightarrow aMb$ $M \to A$ $M \rightarrow B$ $A \rightarrow aA$ $A \rightarrow e$ $B \to bB$ $B \to e$

Equivalency

- A type of automata (as a *recognition* device) can be made to recognize languages generated by a particular grammar (as a *generation* device)
- E.g., the class of languages accepted by finite automata is exactly the class of languages generated by regular grammars



Generalization

- Non-deterministic automata (allows multiple transitions out of a state)
- Pushdown automata (context free languages)
- Stochastic grammar
- The whole area of syntactic pattern recognition
- ♦ etc. etc.



String (Pattern) Grammar

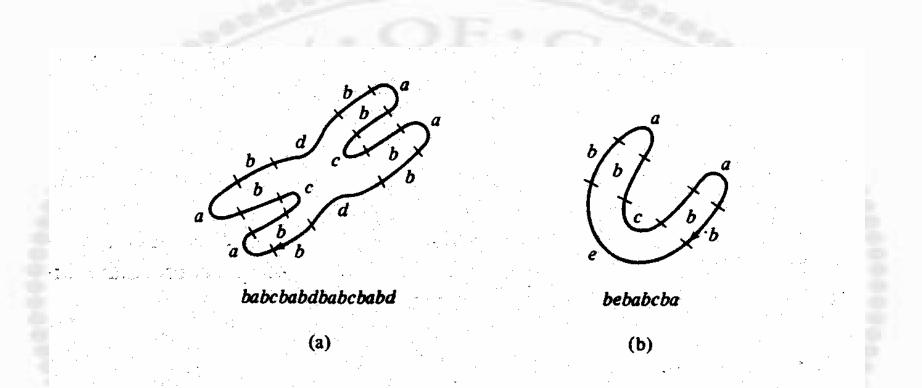


Figure 3.19. (a) Submedium chromosome; (b) telocentric chromosome.



 $G = (V_N, V_T, P, \{\langle \text{submedian chromosome} \rangle, \langle \text{telocentric chromosome} \rangle \})$

where

 $V_N = \{\langle submedian chromosome \rangle, \langle telocentric chromosome \rangle, \langle arm pair \rangle, \langle left part \rangle, \langle right part \rangle, \langle arm \rangle, \langle side \rangle, \langle bottom \rangle \}$

$$V_{T} = \left\{ \bigcap_{a \in C} \left[\left(\bigcup_{a \in C} \left[\left((\bigcup_{a \in C} \left[\left(\bigcup_{a \in C} \left[\left((\bigcup_{a \in C} \left[(\bigcup_{a \in C} \left[$$

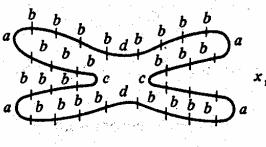
and P:

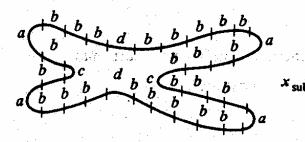
 $\langle submedian chromosome \rangle \longrightarrow \langle arm pair \rangle \langle arm pair \rangle$ $\langle telocentric chromosome \rangle \longrightarrow \langle bottom \rangle \langle arm pair \rangle$ $\langle arm pair \rangle \longrightarrow \langle side \rangle \langle arm pair \rangle$



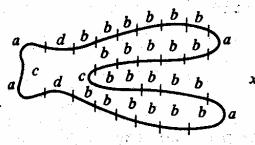
 $\langle \operatorname{arm pair} \rangle \longrightarrow \langle \operatorname{arm pair} \rangle \langle \operatorname{side} \rangle$ $\langle \operatorname{arm} \operatorname{pair} \rangle \longrightarrow \langle \operatorname{arm} \rangle \langle \operatorname{right} \operatorname{part} \rangle$ $\langle \operatorname{arm pair} \rangle \longrightarrow \langle \operatorname{left part} \rangle \langle \operatorname{arm} \rangle$ $\langle \text{left part} \rangle \longrightarrow \langle \text{arm} \rangle c$ $\langle right part \rangle \longrightarrow c \langle arm \rangle$ $\langle bottom \rangle \rightarrow b \langle bottom \rangle$ $\langle bottom \rangle \rightarrow \langle bottom \rangle b$ $\langle bottom \rangle \rightarrow e$ $\langle \text{side} \rangle \longrightarrow b \langle \text{side} \rangle$ $\langle side \rangle \rightarrow \langle side \rangle b$ $\langle side \rangle \longrightarrow b$ $\langle \text{side} \rangle \longrightarrow d$ $\langle \operatorname{arm} \rangle \longrightarrow b \langle \operatorname{arm} \rangle$ $\langle \operatorname{arm} \rangle \longrightarrow \langle \operatorname{arm} \rangle b$ $\langle \operatorname{arm} \rangle \longrightarrow a$







(a)



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Figure 3.20. Three sample chromosomes: (a) median string representation; (b) submedian representation; (c) arcocentric string representation.

(c)

(b)



Tree Grammar

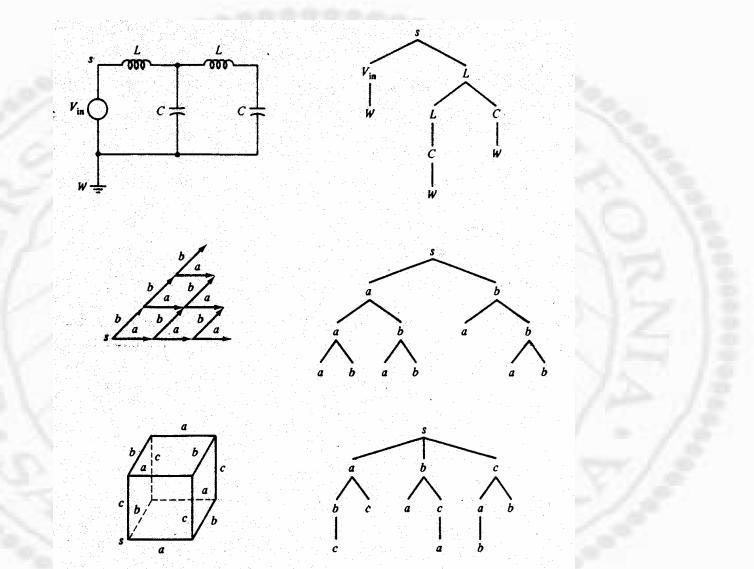


Figure 4.6. (a) Patterns; (b) corresponding tree representations.



Difficulty

- How to extract primitives reliably from images?
- How to parse the the extracted primitives?
 How to correct errors?

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Figure 3.2. Cursive strokes of the word "globe."

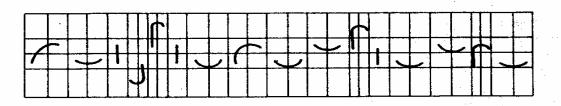


Figure 3.3. Stroke-sequence representation of the word "globe."



Hidden Markov Model

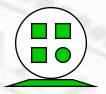
- *Internal states*: short-time, steady, wellbehaved part of a signal or an image
- State transition: characterize how one state evolves into the next
- *Initial states*: the starting point of evolution
 "Hidden" implies that the states are usually not directly observable, but will influence the behavior of the model



An example - urn-and-ball-model



urn1urn2urn3P(R)=.P(R)=.P(R)=.P(B)=.P(B)=.P(B)=.



urnN P(R)=. P(B)=.

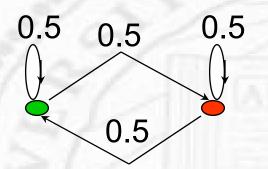
Time123...Turn (hidden) stateu1u9u5...u7ball (observation)RBY...G



Another example - competing HMMs

2-fair coins

2-biased coins



P(H) = 0.5 P(H) = 0.5P(T) = 0.5 P(T) = 0.5 0.5 0.5 0.5 0.5 0.5 0.5 0.5

P(H)=0.8 P(H)=0.2 P(T)=0.2 P(T)=0.8

Observation: HTTHTHH ..

1. Which model best describes the observation?

2. What are the most likely state transitions over time?



Terminology

□ T: length of the observation sequence □ N: number of states (urns) in the model □ M: number of observation symbols (colors) $\{O_1, O_2, ..., O_T\}$ • O: observations (colors) $\{q_1, q_2, ..., q_N\}$ Q: states (urns) $\{v_1, v_2, \dots, v_M\}$ □ V: set of possible symbols $a_{ij} = p(q_j^{t+1}|q_i^t)$ □ A: state transition density B: observation symbol density in a state $b_j(k) = p(v_k^t | q_j^t)$ $\square \pi$: initial state distribution $\pi_i = P(q_i^{-1})$



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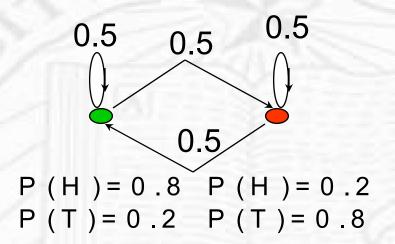
Three Problems of HMMs

Evaluation: Given an observation sequence O, and a model λ = (A, B, π) how to compute P(O|λ) *Estimation*: Given an observation sequence O, and a model λ = (A, B, π) how to choose a state sequence I = i₁, i₂,..., i_T which is optimal *Training*: How to adjust the model parameters to maximize λ = (A, B, π) P(O|λ)



Example - Evaluation

2-biased coins



Observation: H T T H T H H

How good can the states explain the output?

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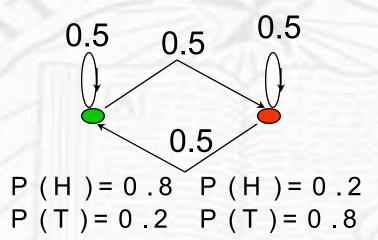
Intuition: Given a Model ...

- A good (or likely) sequence of states is one that
 - Each state can well explain the corresponding observation
 - Each transition has high likelihood to happen
- A bad (or unlikely) sequence of states does not have either (or both) properties
- A probabilistic framework
 - Enumerate all possible state transitions in a model to determine how good is a model



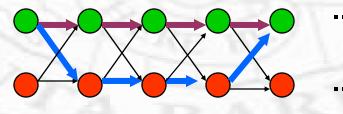
Example - Evaluation

2-biased coins



Observation: H H H H H H H

0.8*0.5*0.8*0.5*0.8*0.5*0.8*0.5*0.8



0.8*0.5*0.2*0.5*0.2*0.5*0.2*0.5*0.8



Evaluation

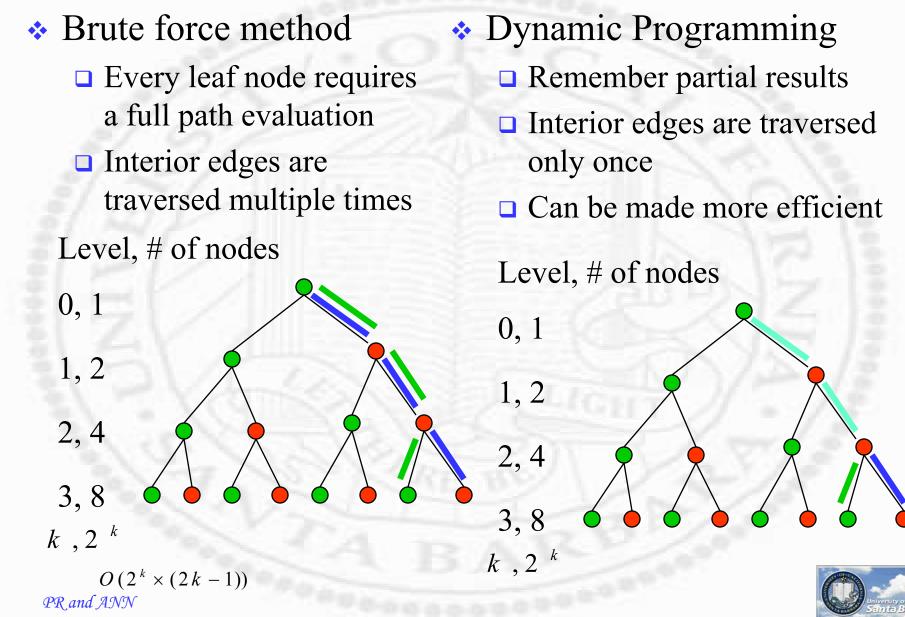
- Enumerates all possible state transitions in T steps
- For each possible state sequence, compute the possibility of observing O

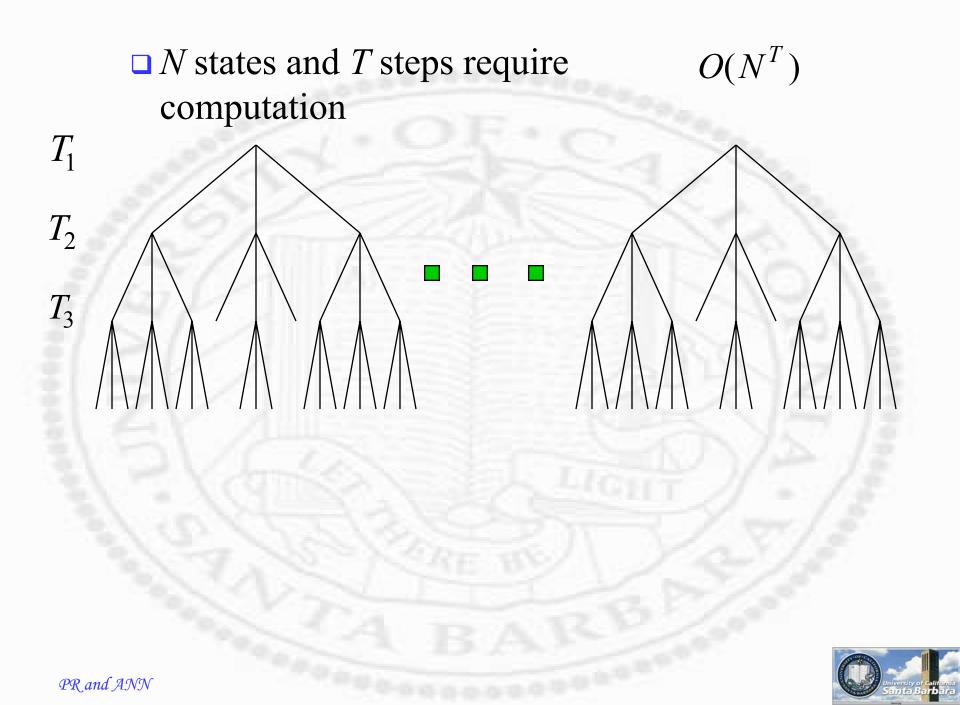
$$P(O \mid \lambda) = \sum_{allI} P(O, I \mid \lambda) = \sum_{allI} p(O \mid I, \lambda) p(I \mid \lambda)$$
$$P(O \mid I, \lambda) = b_{i_1}(O_1) b_{i_2}(O_2) \dots b_{i_T}(O_T)$$

 $p(I \mid \lambda) = \pi_{i_1} a_{i_1 i_2} a_{i_2 i_3} \dots a_{i_{T-1} i_T}$ $P(O \mid \lambda) = \sum_{allI} \pi_{i_1} b_{i_1} (O_1) a_{i_1 i_2} b_{i_2} (O_2) a_{i_2 i_3} \dots a_{i_{T-1} i_T} b_{i_T} (O_T)$

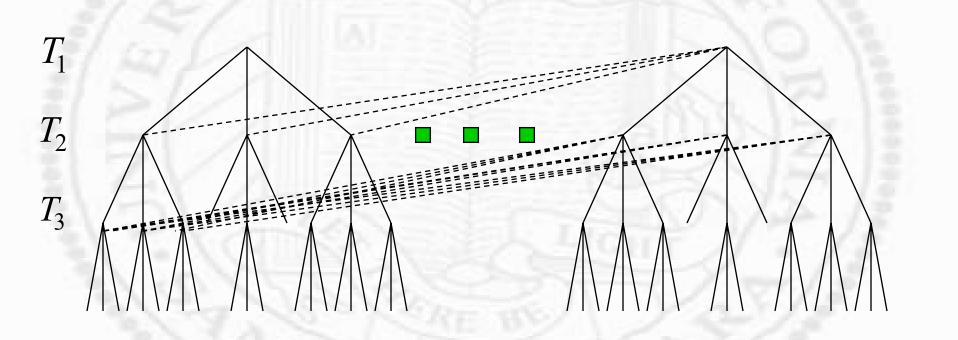


Graphical Interpretation





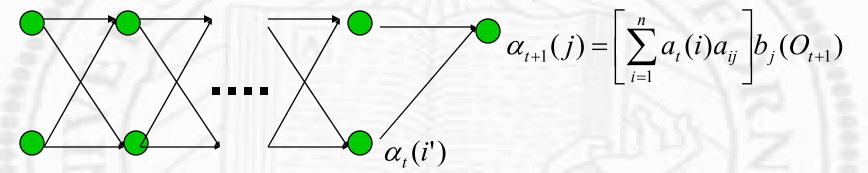
Efficient forward -backward procedure (based on dynamic programming)
 Remember *all* costs





Intuition

* At each node, the cost is the sum of all possible paths from start to that node $\alpha_t(i)$

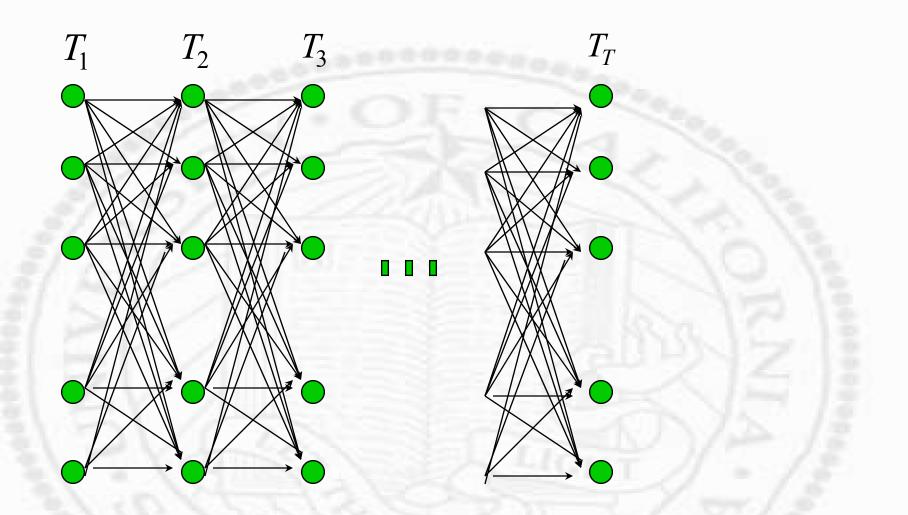


To proceed one more stage (t to t+1), the total cost is

- □ Total cost at stage *t*
- □ Transition cost from *t* to t+1 Extending all paths at once!

Observation cost at stage t+1





N states per stage/Tstages



Forward approach

$$\alpha_{t}(i) = p(O_{1}, O_{2}, ..., O_{t}, i_{t} = q_{i} | \lambda)$$

$$1. \alpha_{1}(i) = \pi_{i} b_{i}(O_{1}) \qquad 1 \le i \le N$$

$$2. \text{ for } t = 1, 2, ..., T - 1 \qquad 1 \le j \le N$$

$$\alpha_{t+1}(j) = \left[\sum_{i=1}^{N} \alpha_{t}(i) a_{ij}\right] b_{j}(O_{t+1})$$

$$3. P(O | \lambda) = \sum_{i=1}^{N} \alpha_{T}(i)$$

* Backward approach

$$\beta_{t}(i) = p(O_{t+1}, O_{t+2}, ..., O_{T} | i_{t} = q_{i}, \lambda)$$

$$1. \beta_{T}(i) = 1 \qquad 1 \le i \le N$$

$$2. \text{ for } t = T - 1, T - 2, ..., 1 \qquad 1 \le i \le N$$

$$\beta_{t}(i) = \sum_{j=1}^{N} a_{ij}b_{j}(O_{t+1})\beta_{t+1}(j)$$

$$3. P(O|\lambda) = \sum_{i=1}^{N} \beta_{1}(i)$$
PR and ANN



Estimation #1

Optimize each stage *individually Does not* take into consideration if a path can be built with those individual states
Usually fine, but *not* if a_{ij} is zero for some transition *Does not* take into consideration the

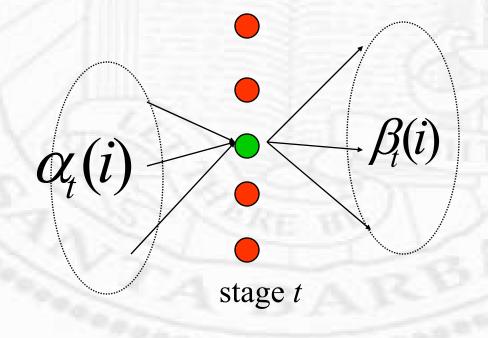
transition costs between those states

A localized, greedy approach



Estimation #1

$$\gamma_{t}(i) = p(q_{t} = i \mid O, \lambda) = \frac{p(O, q_{t} = i, \lambda)}{p(O, \lambda)} = \frac{p(O, q_{t} = i \mid \lambda) p(\lambda)}{p(O \mid \lambda) p(\lambda)} = \frac{p(O, q_{t} = i \mid \lambda)}{p(O \mid \lambda)}$$
$$= \frac{p(O, q_{t} = i \mid \lambda)}{\sum_{j=1}^{N} p(O, q_{t} = j \mid \lambda)} = \frac{\alpha_{t}(i)\beta_{t}(i)}{\sum_{j=1}^{N} \alpha_{t}(j)\beta_{t}(j)}$$
All possible ways through stage t via node i All possible ways through stage t





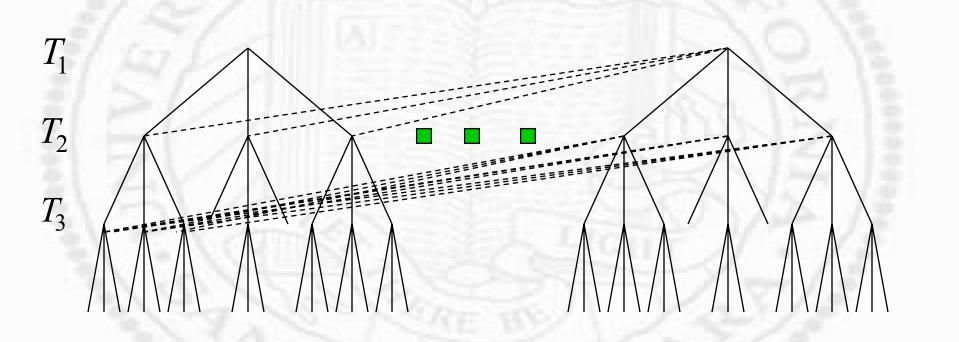
Estimation 2

Find the best state sequence (path)
 Viterbi algorithm: a dynamic programming solution for finding the shortest (best) path

1. $\delta_{1}(i) = \pi_{i}b_{i}(O_{1})$ $\psi_{1}(i) = 0$ 2. $\delta_{t}(j) = \max_{1 \le i \le N} [\delta_{t-1}(i)a_{ij}]b_{j}(O_{t})$ $\psi_{t}(j) = \arg\max_{1 \le i \le N} [\delta_{t-1}(i)a_{ij}]$ 3. $\operatorname{cost} = \max_{1 \le i \le N} [\delta_{T}(i)]$ path through back tracking on ψ

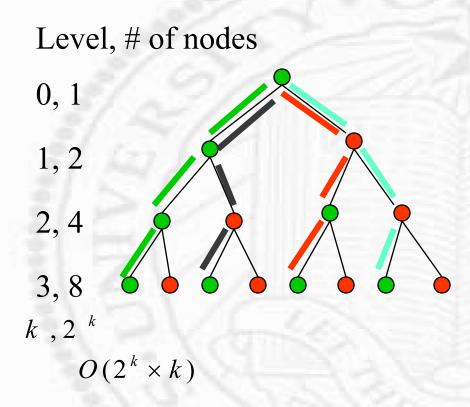


Efficient forward -backward procedure (based on dynamic programming) Remember the *best* cost instead of *all* cost





Graphical Illustration



1,1,1,1 (green) 1,1,2,1 (black) 1,2,1,1 (red) 1,2,2,1 (cyan) All go from 1,x,x,1 Only the best path needs to be remembered

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Training

Observations: All known or partially known
Structures: known or not known
Parameters: known or not known
If we don't know the structure (hence, not parameters either), it is an arbitrarily hard problem

You better have some domain knowledge or some hypothesis of the structures, otherwise, you are doomed



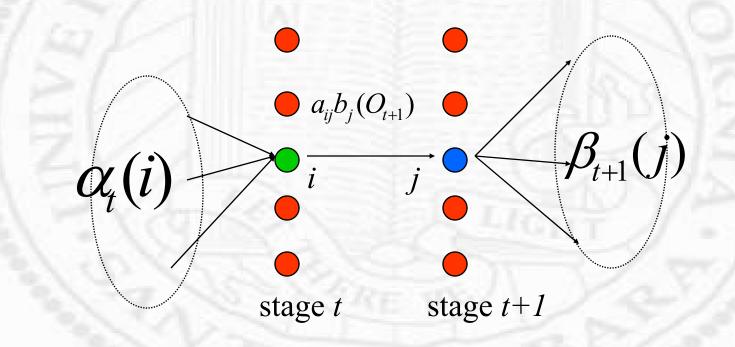
Training

A much more reasonable assumption Observations are available for training Network structure is known Need to determine or fine tune parameters > Initial probability > Transition probability > Observation probability Formulated as EM optimization

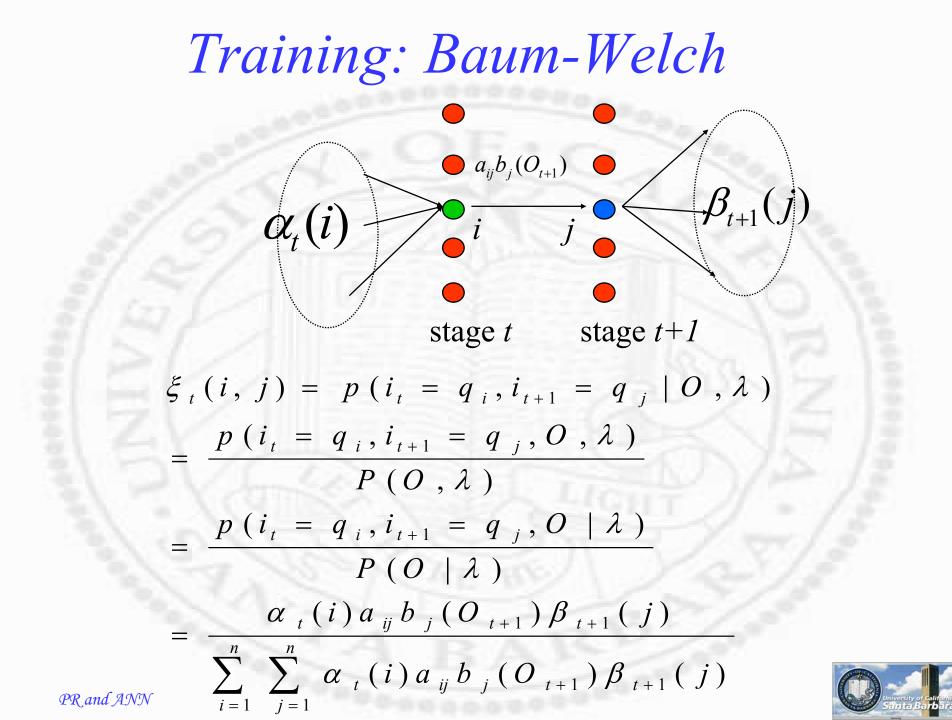


Training: Baum-Welch

Probability of in state *i* at time *t* and state *j* at time *t*+1







More definitions

\$\gam{\gam{t}}{\express{t}}_{i}(i) = \sum_{j=1}^{n} \xi_{t}^{\express{t}}(i,j)\$ Being in state *i* at time *t* \$\sum_{j=1}^{T-1} \gam{t}{\express{t}}_{t}(i)\$ Expected number of transitions made from state \$\vee{q}_{i}\$

$\square \sum_{t=1}^{T-1} \xi_t(i, j) \text{ Expected number of transitions from}$

state q_i to state q_j



Intuition

- * π: expected frequency (# of times) in state *i* at time *t*=1
- * a_{ij}: expected number of transition from *i* to *j*/expected transition from *I*
- *B_j(k)*: expected number of times in state *j* and observing *k*/expected number of times in state *j*



Baum-Welch procedure

□ iterate on $\hat{\lambda} = (\hat{\pi}, \hat{a}_{ij}, \hat{b}_j(k))$ $\hat{\pi}_i$ $\gamma_1(i)$ $\frac{\sum_{t=1}^{T-1} \xi_t(i, j)}{\sum_{t=1}^{T-1} \gamma_t(i)}$ \hat{a}_{ij} $\sum_{t=1}^{T-1} \gamma_t(i)$ $\begin{array}{c} \iota = 1 \\ 0 \\ t = k \\ \hline T \\ \Sigma \\ t = 1 \end{array}^{t} \left(i \right)
 \end{array}$ $\hat{b}_i(k)$



